On a Characteristic of Data Transmission with Three or More Priority Levels in Bus Network

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Abstract: This paper presents a characteristic of data transmission in bus network in which each datum has a priority and the number of priority levels are three or more. In this bus network, each station is connected to a bus through a switching-hub which establishes a link from a sending station to a receiving station. When the receiving station is sending data or receiving data, the link cannot be established and the sending station has to wait a little waiting time (stand-by time). Then, by defining the link probability as a constant, the bus network is analyzed.

Taking note of one of the sending stations with $r (r \geq 3)$ sending buffers for data with priority, in steady state, a balance equation of the probabilities of the number of waiting data at the ready time of sending is expressed. From the equation, for each priority, probabilities of the number of waiting data, mean of the number of waiting data in the sending buffer and mean waiting time which is the time from the arrival time of a datum at the sending buffer to the beginning time of sending the datum are derived. The calculated values of mean waiting time are illustrated and mean waiting times with priority processing are compared with mean waiting time without priority processing. As a result, the influence of the link probability on data transmission with priority processing is clarified.

Key words: Bus network, priority processing, probability of the number of waiting data, mean of the number of waiting data, mean waiting time.

1. Introduction

In bus network which is a primary topology for LAN (Local Area Network), recently, stations are connected to a bus through a switching hub which establishes a link from a sending station to a receiving station. As studies in bus network, the proposal of the dynamic bandwidth allocation [1] for Ethernet PON(Passive Optical Network), the time synchronization on Gb/s Ethernet [2] which is necessary for the wide distributed system and the host configuration mechanism to network in Ethernet and wireless LAN [3] are studied. For application of Ethernet, an application method in industry [4], the performance of industrial Ethernet [5] and the performance of TDMA based network [6] are reported. And, about ad hoc network, impact of mobility [7], [8] are studied.

In bus network with a switching hub, a sending station has to wait for sending when the receiving station is sending or receiving data. I already reported a characteristic of data transmission in bus network under the condition that the link probability is defined as a constant [9], [10].

In the case of transmission of different kinds of data, such as text, sound and graphics, priority processing is necessary. For priority processing, the proposal of MAC protocol with priority adjustment [11] and the simulation of MIN with priority routing [12] are studied. And, CAN with priority [13] and ad hoc network
with priority [14] are reported. I also reported bus network with two priority levels [15]. However, in order to consider a characteristic of bus network with more priority levels, the analysis in the case of two priority levels is not sufficient.

In this paper, a characteristic of data transmission with priority processing in bus network which has three or more priority levels is clarified. Taking note of one of the sending stations which has $r(r\geq3)$ sending buffers for data with priority, probabilities of the number of waiting data, mean of the number of waiting data in the sending buffer and mean waiting time from the arrival time at the sending buffer to the beginning time of sending are derived. Furthermore, the calculated values of mean waiting time are illustrated and mean waiting times with priority processing are compared with mean waiting time without priority processing. As a result, the influence of the link probability on data transmission with priority processing is clarified.

2. Modeling Of Bus Network with Three or More Priority Levels

The bus network shown in Fig. 1 is analyzed.

![Fig. 1. Bus network with three or more priority levels.](image)

In the network, each station has $r(r\geq3)$ infinite buffers for priority level data, which are called as sending buffers. Each priority level data arrive at a sending buffer corresponding to priority level in different
Poisson process. Priority level 1 is the highest level and priority level \( r \) is the lowest. The arrival ratio for priority level \( i \) \((i=1,2,\ldots,r)\) data is \( \lambda_i \). At the ready time of sending, if priority level 1 data exist in the sending buffer and a link to a receiving station is established, the station sends the first priority level 1 datum in the sending buffer. If not established, the station has a little waiting time (called as stand-by time). If there is no priority level 1 data, the station checks priority level 2 data. If priority level 2 data exist and the link is established, the station sends the first priority level 2 datum. If not established, the station has stand-by time. If there is no priority level 2 data, the station checks priority level 3 data. The same processing is performed, and finally, if priority level \( r \) data exist and the link is established, the station sends the first priority level \( r \) datum. If not established, the station has stand-by time. Defining the link probability as a constant \( q \), the flowchart of data transmission is shown in Fig. 2.

Taking note of one of the sending stations, the sending station is modeled as Fig. 3.

Transmission time is arbitrary distributed and the distribution function (DF) of the transmission time for each priority level data is denoted by \( H_i(t) \) \((i=1,2,\ldots,r)\). The stand-by time is arbitrary distributed and the DF of the stand-by time is denoted by \( H_c(t) \).

### 3. Probability of the Number of Waiting Data

#### 3.1. Balance Equation

Let \( P_{n_1,n_2,\ldots,n_r} \) be the probability that the number of priority level \( i \) \((i=1,2,\ldots,r)\) waiting data in the sending buffers is \( n_i \) at the ready time of sending. In steady state, the following balance equation is obtained.

\[
P_{n_1,n_2,\ldots,n_r} = q \sum_{n_r=0}^{n_r'} a_r(n_r') \sum_{n_{r-1}=0}^{n_{r-1}'} a_{r-1}(n_{r-1}') \prod_{i=r-1}^{r-2} \sum_{n_i=0}^{n_i'} a_i(n_i') \prod_{j=i}^{1} \sum_{n_j=0}^{n_j'} \lambda_j(n_j'-n_j,t) P_{n_1,n_2,\ldots,n_r} dH_i(t)
\]

\[
+ q \sum_{n_r=0}^{n_r'} a_r(n_r') \sum_{n_{r-1}=0}^{n_{r-1}'} a_{r-1}(n_{r-1}') \prod_{i=r-1}^{r-2} \sum_{n_i=0}^{n_i'} a_i(n_i') \prod_{j=i}^{1} \sum_{n_j=0}^{n_j'} \lambda_j(n_j'-n_j,t) P_{n_1,n_2,\ldots,n_r} dH_i(t)
\]

\[
+ \cdots
\]

\[
+ q \sum_{n_r=0}^{n_r'} a_r(n_r') \sum_{n_{r-1}=0}^{n_{r-1}'} a_{r-1}(n_{r-1}') \prod_{i=r-1}^{r-2} \sum_{n_i=0}^{n_i'} a_i(n_i') \prod_{j=i}^{1} \sum_{n_j=0}^{n_j'} \lambda_j(n_j'-n_j,t) P_{n_1,n_2,\ldots,n_r} dH_i(t)
\]

\[
+ (1-q) \sum_{n_r=0}^{n_r'} a_r(n_r') \sum_{n_{r-1}=0}^{n_{r-1}'} a_{r-1}(n_{r-1}') \prod_{i=r-1}^{r-2} \sum_{n_i=0}^{n_i'} a_i(n_i') \prod_{j=i}^{1} \sum_{n_j=0}^{n_j'} \lambda_j(n_j'-n_j,t) P_{n_1,n_2,\ldots,n_r} dH_i(t)
\]

(1)
In equation (1), \( \alpha_i(n,t) \) is the probability that \( n \) priority level \( i \) data arrive at the sending buffer within time \( t \)

\[
\alpha_i(n,t) = \frac{e^{-\lambda_i t} (\lambda_i t)^n}{n!} \quad (i=1,2,\ldots,r)
\]  

(2)

Here, in order to make analytical procedure plain, let the number of priority levels be 3, i.e., \( r=3 \).

Taking \( \sum_{n_1=0}^{\infty} \sum_{n_2=0}^{\infty} \sum_{n_3=0}^{\infty} \) after multiplying \( x_1^{n_1}x_2^{n_2}x_3^{n_3} \) in both sides of equation (1) under \( r=3 \) and transforming into a generating function form, the following equation is obtained.

\[
G_1(x_1,x_2,x_3) = \frac{1}{1 - \frac{q}{x_1} H_1^*(z) - (1-q) H_c^*(z)} \times \left[ G_2(x_2,x_3) \left\{ \frac{q}{x_2} H_2^*(z) + (1-q) H_c^*(z) - 1 \right\} 
+ G_3(x_3) \left\{ \frac{q}{x_3} H_3^*(z) + (1-q) H_c^*(z) - 1 \right\} + P_{0,0,0} \left\{ H_c^*(z) - 1 \right\} \right]
\]  

(3)

where \( H_i^*(z) \) (i=1,2,3) and \( H_c^*(z) \) are the Laplace-Stieltjes Transformation(LST) of \( H_i(t) \) and \( H_c(t) \), respectively, and \( z = \lambda_1(1-x_1) + \lambda_2(1-x_2) + \lambda_3(1-x_3) \).

The generating functions of the probabilities of the number of waiting data are defined as follows.

\[
G_1(x_1,x_2,x_3) = \sum_{n_1=1}^{\infty} \sum_{n_2=0}^{\infty} \sum_{n_3=0}^{\infty} P_{n_1,n_2,n_3} x_1^{n_1} x_2^{n_2} x_3^{n_3}
\]  

(4a)

\[
G_2(x_2,x_3) = \sum_{n_2=1}^{\infty} \sum_{n_3=0}^{\infty} P_{0,n_2,n_3} x_2^{n_2} x_3^{n_3}
\]  

(4b)

\[
G_3(x_3) = \sum_{n_3=1}^{\infty} P_{0,0,n_3} x_3^{n_3}
\]  

(4c)

3.2. Probabilities of the Number of Waiting Data

When \( x_1 = x_2 = x_3 = 1 \) in equation (3), the following equations are obtained.

\[
\left\{ q - q \lambda_1 h_1 - (1-q) \lambda_1 h_c \right\} G_1(1,1,1) - \left\{ q \lambda_1 h_2 + (1-q) \lambda_1 h_c \right\} G_2(1,1) = 0
\]

(5a)

\[
\left\{ q \lambda_2 h_1 + (1-q) \lambda_2 h_c \right\} G_1(1,1,1) - \left\{ q - q \lambda_2 h_2 - (1-q) \lambda_2 h_c \right\} G_2(1,1) + \left\{ q \lambda_2 h_3 + (1-q) \lambda_2 h_c \right\} G_3(1,1) = 0
\]

(5b)
\[
\left\{ q\lambda_3 h_j + (1-q)\lambda_3 h_k \right\} G_1 (1,1,1) + \left\{ q\lambda_3 h_2 + (1-q)\lambda_3 h_k \right\} G_2 (1,1) \\
- \left\{ q - q\lambda_3 h_j - (1-q)\lambda_3 h_k \right\} G_3 (1) + \lambda_3 h_k P_{0,0,0} = 0
\] (5c)

In equation (5a), (5b) and (5c), \( h_i \) and \( h_c \) are means of \( H_i(t) \) and \( H_c(t) \), respectively. Because the sum of all probabilities is equal to 1, equation (6) is obtained.

\[
G_1(1,1,1) + G_2(1,1) + G_3(1) + P_{0,0,0} = 1
\] (6)

From equation (5a), (5b), (5c) and (6), the following probabilities are derived.

\[
G_1(1,1,1) = \frac{\lambda_1 h_c}{q\{1-(\lambda_1 h_j + \lambda_2 h_j + \lambda_3 h_j) + (\lambda_1 + \lambda_2 + \lambda_3)h_k\}}
\] (7a)

\[
G_2(1,1) = \frac{\lambda_2 h_c}{q\{1-(\lambda_1 h_j + \lambda_2 h_j + \lambda_3 h_j) + (\lambda_1 + \lambda_2 + \lambda_3)h_k\}}
\] (7b)

\[
G_3(1) = \frac{\lambda_3 h_c}{q\{1-(\lambda_1 h_j + \lambda_2 h_j + \lambda_3 h_j) + (\lambda_1 + \lambda_2 + \lambda_3)h_k\}}
\] (7c)

\[
P_{0,0,0} = 1 - G_1(1,1,1) - G_2(1,1) - G_3(1)
\] (7d)

\( G_1(1,1,1) \) is the probability that the number of priority level 1 waiting data is one or more at the ready time of sending. \( G_2(1,1) \) is the probability that the number of priority level 1 waiting data is 0 and priority level 2 data is one or more. \( G_3(1) \) is the probability that the number of priority level 1 and level 2 waiting data is 0 and priority level 3 data is one or more. \( P_{0,0,0} \) is the probability that there is no waiting data.

4. Mean of the Number of Waiting Data

4.1. Priority Level 1

From equation (3), \( \frac{d}{dx_i} G_i(x_i,1,1)|_{x_i=1} \) which is mean of the number of priority level 1 waiting data in the sending buffer can be derived.

From equation (3), \( G_i(x_i,1,1) \) is expressed by equation (8), (9a) and (9b).

\[
G_i(x_i,1,1) = \frac{A_i(x_i)}{B_i(x_i)}
\] (8)

\[
A_i(x_i) = G_2(1,1) \left\{ qH_2^*(z_j) + (1-q)H_c^*(z_j) - 1 \right\} \\
+ G_3(1) \left\{ qH_3^*(z_j) + (1-q)H_c^*(z_j) - 1 \right\} \\
+ P_{0,0,0} \left\{ H_c^*(z_j) - 1 \right\}
\] (9a)
where \( z_j = \lambda_j (1 - x_j) \).

When \( x_j = 1 \) after differentiating equation (8), the following equation is obtained

\[
\frac{d}{dx_j} G_j(x_j, 1) \bigg|_{x_j=1} = \frac{A_i''(1) B_i'(1) - A_i'(1) B_i''(1)}{2 \{B_i'(1)\}^2} \tag{10}
\]

where

\[
A_i'(1) = G_1(1, 1) \{ q \lambda_i h_j + (1 - q) \lambda_i h_c \} + G_j(1) \{ q \lambda_i h_j + (1 - q) \lambda_i h_c \} + P_{0,0,0} \lambda_i h_c
\]

\[\text{and} \]

\[
A_i''(1) = G_2(1, 1) \{ q \lambda_i^2 h_j^{(2)} + (1 - q) \lambda_i^2 h_c^{(2)} \} + G_j(1) \{ q \lambda_i^2 h_j^{(2)} + (1 - q) \lambda_i^2 h_c^{(2)} \} + P_{0,0,0} \lambda_i^2 h_c^{(2)}
\]

\[\text{and} \]

\[
B_i'(1) = q - q \lambda_i h_j - (1 - q) \lambda_i h_c
\]

\[\text{and} \]

\[
B_i''(1) = -2q + 2q \lambda_i h_j - q \lambda_i^2 h_j^{(2)} - (1 - q) \lambda_i^2 h_c^{(2)}
\]

\( h_j^{(2)} \) and \( h_c^{(2)} \) are the second moments of the DFs of \( H_i(t) \) and \( H_c(t) \), respectively.

### 4.2. Priority Level 2

Using \( x_j = \delta_i(x_2, x_3) \) which satisfies that the denominator of equation (3) is equal to 0, the following equation is obtained.

\[
G_2(x_2, x_3) = \frac{G_j(x_3) \{ \frac{1}{x_3} q H_j^*(z) + (1 - q) H_c^*(z) - 1 \} + P_{0,0,0} \{ H_{0,0}^*(z) - 1 \}}{1 - \frac{q}{x_2} H_j^*(z) - (1 - q) H_c^*(z)} \tag{13}
\]

where \( z = \lambda_1 \{ 1 - \delta_i(x_2, x_3) \} + \lambda_2 (1 - x_2) + \lambda_3 (1 - x_3) \).

and \( \delta_i(x_2, x_3) \) satisfies the following equation.

\[
1 - \frac{q}{\delta_i(x_2, x_3)} H_i^*(z) - (1 - q) H_c^*(z) = 0
\]

\[\text{where} \]

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From equation (13), \( \frac{d}{dx_2} G_i(x_2,1) \) which is mean of the number of priority level 2 waiting data in the case that the number of priority level 1 waiting data is 0 can be derived. From equation (13), \( G_i(x_2,1) \) is expressed by equation (15), (16a) and (16b).

\[
G_i(x_2,1) = \frac{A_2(x_2)}{B_2(x_2)}
\]

(15)

\[
A_2(x_2) = G_3(1)\{qH'_c(z_2) + (1 - q)H'_e(z_2) - 1\} + P_{0,0,0}\{H'_c(z_2) - 1\}
\]

(16a)

\[
B_2(x_2) = 1 - \frac{q}{x_2}H'_c(z_2) - (1 - q)H'_e(z_2)
\]

(16b)

where \( z_2 = \lambda_1 \{1 - \delta(x_2,1)\} + \lambda_2 (1 - x_2) \).

When \( x_2 = 1 \) after differentiating equation (15), the following equation is obtained.

\[
\frac{d}{dx_2} G_i(x_2,1) \bigg|_{x_2=1} = \frac{A_2''(1)B_2'(1) - A_2'(1)B_2''(1)}{2\{B_2'(1)\}^2}
\]

(17)

\[
A_2'(1) = G_3(1)\{-q\tilde{z}_{20}h_3 - (1 - q)\tilde{z}_{20}h_e\} - P_{0,0,0}\tilde{z}_{20}h_c
\]

(18a)

\[
B_2'(1) = q + q\tilde{z}_{20}h_2 + (1 - q)\tilde{z}_{20}h_c
\]

(18b)

where \( \tilde{z}_{20} = \tilde{z}_2 \bigg|_{x_2=1} \)

and

\[
A_2''(1) = G_3(1)\{q(\tilde{z}_{20})^2h_3^{(2)} - q\tilde{z}_{20}h_3 + (1 - q)(\tilde{z}_{20})^2h_c^{(2)}
\]

\[
-(1 - q)\tilde{z}_{20}h_e + P_{0,0,0}\{(\tilde{z}_{20})^2h_c^{(2)} - \tilde{z}_{20}h_c\}
\]

(19a)

\[
B_2''(1) = -2q - 2q\tilde{z}_{20}h_2 - q(\tilde{z}_{20})^2h_2^{(2)} + q\tilde{z}_{20}h_2
\]

\[
-(1 - q)(\tilde{z}_{20})^2h_c^{(2)} + (1 - q)\tilde{z}_{20}h_c
\]

(19b)

where \( \tilde{z}_{20} = \tilde{z}_2 \bigg|_{x_2=1} \).

On the other hand, from \( \tilde{z}_2 = \lambda_1 \{1 - \delta(x_2)\} + \lambda_2 (1 - x_2) \).

and equation (14) under \( x_3 = 1 \), the following equations are obtained.

\[
\delta'_{10} = \frac{q\lambda_2h_1 + (1 - q)\lambda_2h_c}{q - q\lambda_2h_1 - (1 - q)\lambda_2h_c}
\]

(20a)
From these equations, $\tilde{z}''_{20}$ and $\tilde{z}'_{20}$ can be calculated.

**4.3. Priority Level 3**

Using $x_3 = \delta_3(x_j)$ which satisfies that the denominator of equation (13) is equal to 0, the following equation is obtained.

$$G_j(x_3) = \frac{P_{0,0,0} \left\{ H_c^*\left( \bar{z} \right) - 1 \right\}}{1 - \frac{x_3 \cdot q}{H_j^*\left( \bar{z} \right)} - (1 - q) \cdot H_c^*\left( \bar{z} \right)}$$

(21)

where $\bar{z} = \lambda_1 \left\{ 1 - \delta_1 (\delta_2 (x_j), x_j) \right\} + \lambda_2 \left\{ 1 - \delta_2 (x_j) \right\} + \lambda_3 \left\{ 1 - x_j \right\}$.

and $\delta_1 (\delta_2 (x_j), x_j)$ and $\delta_2 (x_j)$ satisfy the following equations.

$$1 - \frac{q}{\delta_1 (\delta_2 (x_j), x_j)} \cdot H_j^* \left( \bar{z} \right) - (1 - q) \cdot H_c^* \left( \bar{z} \right) = 0$$

(22)

$$1 - \frac{q}{\delta_2 (x_j)} \cdot H_j^* \left( \bar{z} \right) - (1 - q) \cdot H_c^* \left( \bar{z} \right) = 0$$

(23)

From equation (21), $\frac{d}{dx_j} G_j(x_3) \big|_{x_j=1}$ which is mean of the number of priority level 2 waiting data in the case that the number of priority level 1 and level 2 waiting data is 0 can be derived. From equation (21), $G_j(x_3)$ is expressed by equation (24), (25a) and (25b).

$$G_j(x_3) = \frac{A_3(x_j)}{B_3(x_3)}$$

(24)

$$A_3(x_j) = P_{0,0,0} \left\{ H_c^*\left( \bar{z} \right) - 1 \right\}$$

(25a)

$$B_3(x_3) = 1 - \frac{x_3 \cdot q}{H_j^*\left( \bar{z} \right)} - (1 - q) \cdot H_c^*\left( \bar{z} \right)$$

(25b)
When \( x_3 = 1 \) after differentiating equation (21), the following equation is obtained.

\[
\frac{d}{dx_3} G_3(x_3) \bigg|_{x_3=1} = \frac{A_3''(1)B_3'(1)-A_3'(1)B_3''(1)}{2\{B_3'(1)\}^2}
\]  
\( (26) \)

\[
A_3'(1) = -P_{000} \tilde{z}_0' h_c
\]  
\( (27a) \)

\[
B_3'(1) = q + q \tilde{z}_0'h_3 + (1-q)\tilde{z}_0'h_c
\]  
\( (27b) \)

where \( \tilde{z}_0' = \tilde{z}_0 \bigg|_{x_3=1} \).

and

\[
A_3''(1) = P_{000} \{ (\tilde{z}_0')^2 h_c^{(2)} - \tilde{z}_0' h_c \}
\]  
\( (28a) \)

\[
B_3''(1) = -2q - 2q \tilde{z}_0'h_3 - q(\tilde{z}_0')^2 h_3^{(2)} + q \tilde{z}_0'h_3
\]

\[
- (1-q)(\tilde{z}_0')^2 h_c^{(2)} + (1-q)\tilde{z}_0'h_c
\]  
\( (28b) \)

where \( \tilde{z}_0'' = \tilde{z}_0' \bigg|_{x_3=1} \).

On the other hand,

\[ z = \lambda_1 \{ 1 - \delta_1(\delta_2(x_3),x_3) \} + \lambda_2 \{ 1 - \delta_2(x_3) \} + \lambda_3 (1-x_3) \]

and equation (22), (23), \( \delta''_{20}, \delta''_{130}, \tilde{z}_0' \) and \( \delta''_{20}, \delta''_{130}, \tilde{z}_0'' \) can be derived.

Here, \( \delta''_{20} = \delta''_{2}(x_3) \bigg|_{x_3=1} \), \( \delta''_{130} = \delta''_{1}(\delta_2(x_3),x_3) \bigg|_{x_3=1} \), \( \tilde{z}_0' = \tilde{z}_0' \bigg|_{x_3=1} \),

\[ \delta''_{20} = \delta''_{2}(x_3) \bigg|_{x_3=1} \), \( \delta''_{130} = \delta''_{1}(\delta_2(x_3),x_3) \bigg|_{x_3=1} \), \( \tilde{z}_0'' = \tilde{z}_0'' \bigg|_{x_3=1} \).

5. Mean Waiting Time

5.1. Priority Level 1

The probability that the number of priority level 1 waiting data at the end of sending data is \( n_1' \) is denoted by \( R_{1n_1'} \), and then the following equation is obtained.

\[
R_{1n_1'} = \frac{1}{G_i(1,1,1)} \left[ \sum_{n_1=1}^{\infty} a_1(n_1'-n_1+1,t) \sum_{n_2=0}^{\infty} \sum_{n_3=0}^{\infty} P_{n_1,n_2,n_3} dH_i(t) \right]
\]  
\( (29) \)
Taking $\sum_{n_i=0}^{\infty}$ after multiplying $x_i^{n_i}$ in both sides in equation (29), and transforming into equation of the GF, equation (30) is obtained.

$$G_{R_i}(x_i) = \frac{1}{x_i G_i(1,1,1)} G_i(x_i,1,1) H_i^*(z_i)$$  \hspace{1cm} (30)

where $G_{R_i}(x_i) = \sum_{n_i=0}^{\infty} R_{in_i} x_i^{n_i}$.

On the other hand, the probability that the number of priority level 1 arrival data during waiting time and sending time for priority level 1 data is $n_i'$ is equal to $R_{in_i}$, so that the following equation is obtained.

$$R_{in_i} = \int_{0}^{\infty} \alpha_i(n_i',t) dW_i(t) \otimes H_i(t)$$  \hspace{1cm} (31)

where $W_i(t)$ is the DF of waiting time for priority level 1 data and $\otimes$ is a convolution operator.

Taking $\sum_{n_i=0}^{\infty}$ after multiplying $x_i^{n_i}$ in both sides of equation (31) and transforming into equation of the GF, equation (32) is derived.

$$G_{R_i}(x_i) = W_i^*(z_i) H_i^*(z_i)$$  \hspace{1cm} (32)

where $W_i^*(z_i)$ is the LST of $W_i(t)$.

From equation (30) and (32), $W_i^*(z)$ is obtained as follows.

$$W_i^*(z_i) = \frac{1}{x_i G_i(1,1,1)} G_i(x_i,1,1)$$  \hspace{1cm} (33)

When $z_i = 0$ after differentiating equation (33) by $z_i$, the left side of equation is obtained as follows.

$$\frac{dW_i^*(z_i)}{dz_i} \bigg|_{z_i=0} = -w_j$$  \hspace{1cm} (34a)

where $w_j$ is mean of $W_i(t)$.

On the other hand, the right side is obtained as follows.

$$\frac{d}{dz_i} \left. \frac{G_i(x_i,1,1)}{x_i G_i(1,1,1)} \right|_{z_i=0} = -\frac{1}{\lambda_i G_i(1,1,1)} \left\{ \frac{d}{dx_j} G_j(x_j,1,1) \bigg|_{x_j=1} - G_j(1,1,1) \right\}$$  \hspace{1cm} (34b)
From equation (34a) and (34b), mean waiting time for priority level 1 data \( w_1 \) is obtained.

\[
w_1 = \frac{1}{\lambda_1 G_1(1,1,1)} \left\{ \frac{d}{dx_1} G_1(x_1,1,1)|_{x_1=1} - G_1(1,1,1) \right\}
\]  

(35)

\( G_1(1,1,1) \) and \( \frac{d}{dx_1} G_1(x_1,1,1)|_{x_1=1} \) have been already derived, so that mean waiting time \( w_1 \) can be calculated.

### 5.2. Priority Level 2 and Priority Level 3

Similar to priority level 1, mean waiting times for priority level 2 and level 3 are obtained as follows.

\[
w_2 = \frac{1}{\lambda_2 G_2(1,1)} \left\{ \frac{d}{dx_2} G_2(x_2,1)|_{x_2=1} - G_2(1,1) \right\}
\]  

(36)

\[
w_3 = \frac{1}{\lambda_3 G_3(1)} \left\{ \frac{d}{dx_3} G_3(x_3)|_{x_3=1} - G_3(1) \right\}
\]  

(37)

\( G_2(1,1), \frac{d}{dx_2} G_2(x_2,1)|_{x_2=1}, G_3(1) \) and \( \frac{d}{dx_3} G_3(x_3)|_{x_3=1} \) have been already derived, so that mean waiting times \( w_2 \) and \( w_3 \) can be calculated.

### 5.3. Numerical Examples

Fig. 4 shows a relation between arrival ratio \( \lambda_1 \) and mean waiting times \( w_1, w_2 \) and \( w_3 \) in the case of \( \lambda_1 = \lambda_2 = \lambda_3 \), \( q = 0.1, 0.5, 0.9 \), \( h_1 = h_2 = h_3 = 5 \), \( h_\text{c} = 1 \) and \( H_i(t) \) and \( H_c(t) \) obey uniform distributions.

![Fig. 4. Mean waiting time](image)

Fig. 4. Mean waiting time \( w_1, w_2 \) and \( w_3 \) (\( h_1 = h_2 = h_3 = 5 \), \( h_\text{c} = 1 \), \( \lambda_1 = \lambda_2 = \lambda_3 \), uniform distribution).
Fig. 5. Mean waiting time $w_1$, $w_2$ and $w_3$ ($h_1 = h_2 = h_3 = 5$, $k_1 = 1$, $\lambda_1 = \lambda_2 = \lambda_3$, exponential distribution).

In Fig. 4, solid lines, dashed lines and chain lines are mean waiting times in the case of $q=0.1$, $0.5$ and $0.9$, respectively. In the same lines, the lower line is mean waiting time $w_1$, the middle line is $w_2$ and the upper line is $w_3$. It is confirmed that the mean waiting time becomes long as the link probability $q$ becomes low and as the arrival ratio $\lambda_i$ becomes high. In the case of $q=0.1$, mean waiting time $w_3$ becomes infinite when $\lambda_1$ is over 0.023. This means that when $q=0.1$ and $\lambda_1$ is over 0.023, the sending buffer for priority level 3 is overflow because there are more arrival data than sending data.

In the case that $H_i(t)$ and $H_c(t)$ obey exponential distributions, mean waiting times are shown in Fig. 5. Comparing Fig. 5 to Fig. 4, mean waiting times become longer, but Fig. 5 shows similar tendency.

Next, a comparison with mean waiting times in the case of $q=0.5$, $\lambda_1 = \lambda_2 = \lambda_3$ and mean waiting time without priority processing is shown in Fig. 6.

Fig. 6. Comparison with the mean waiting times and the mean waiting time without priority processing.
In Fig. 6, bold solid line is mean waiting time without priority processing and thin solid line, dashed line and chain line are mean waiting time $w_1$, $w_2$ and $w_3$ with priority processing in the case of $q=0.5$, $h_j = h_1 = h_3 = 5$, $h_2 = 1$, $\lambda_1 = \lambda_2 = \lambda_3$.

It is confirmed that mean waiting time $w_1$, $w_2$ are shorter, but mean waiting time $w_3$ is longer than mean waiting time without priority processing. As arrival ratio becomes high, differences of mean waiting times with priority processing and mean waiting time without priority processing becomes long. On the basis of mean waiting time without priority processing, the ratios of difference of mean waiting times are -44.1% in $w_1$, -22.1% in $w_2$ and 66.1% in $w_3$ on the average.

6. Conclusion

In this paper, a characteristic of data transmission with three or more priority levels in bus network is clarified. In the analysis, taking note of one station, the link probability that the link from the station to a receiving station can be established is defined as a constant. And then, for each priority, the probabilities of the number of waiting data, the mean of the number of waiting data and the mean waiting times are derived. For the analysis of lower priority levels, roots of the denominator of generating functions are used.

In the analysis, at first, the balance equation with arbitrary priority levels is expressed, but the analysis is performed for three priority levels. As for the analysis for more than three priority levels, the analytical procedure of this paper can be used. But, the number of simultaneous equations increases as the number of priority levels increases.

According to the analysis result, some numerical examples are illustrated, and then the influences of the link probability and the arrival ratio of data on the characteristics of data transmission are clarified. Furthermore, the effect of priority processing is confirmed by comparing with mean waiting time with priority processing and that without priority processing.

The characteristic of data transmission with priority processing in this paper is a base of considering efficiency of bus network.

The following considerations remain for further study:

1. Analyzing the bus network with other priority processing.
2. Clarifying a relation between the link probability and the buffer size of a receiving station.

References


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